

Presheaves on midsymmetrical quantaloids

Remigijus Petras GYLYS (MII)

1. Introduction

In previous paper [2] we studied presheaves and sheaves on an arbitrary quantaloid (category enriched in the category of complete sup-lattices subject to certain laws), where we had choosed the C.J. Mulvey and M. Nawaz theory [3] as a guideline for futher generalizations taking our inspiration in the matrix approach of G. Van den Bossche [1]. In this note we continue those investigations restricting our interest to a special subclass of quantaloids, "midsymmetrical" quantaloids. In this setting it is possible to express the structural matrix of the underlying Q -set of a regular and separated presheaf in terms of its "restriction" and "diagonal" (Theorem 3.8). We are going to detail it in a subsequent paper.

2. Quantaloids and matrices

We first review some of the basic notions.

DEFINITION 2.1. A *quantaloid* is a locally small category Q such that:

- (i) for all u, v objects in Q , the hom-set $Q(u, v)$ is a complete lattice,
- (ii) composition of morphisms of Q (in this paper denoted by $\&$) preserves arbitrary joins in both variables:

$$p \& \bigvee_i q_i = \bigvee_i p \& q_i \text{ and } (\bigvee_i p_i) \& q = \bigvee_i p_i \& q$$

for all morphisms p, q of Q and for all families $(p_i), (q_i)$ of morphisms of Q (forming respective composable pairs). A quantaloid Q will be called *midsymmetrical* whenever it satisfies

- (iii) $p \& (q \& r) \& p' = p \& (r \& q) \& p'$ for all $p \in Q(u, v), q, r \in Q(v, v), p' \in Q(v, v')$.

Note that we use the unconventional left-to-right direction for composition of morphisms. An important example of a one-object midsymmetrical quantaloid is given by the lattice of all closed left-sided (or right-sided) ideals of a non-commutative C^* -algebra.

From now Q will be an arbitrary midsymmetrical quantaloid having a small set of objects. Let Q_0 denote this set and Q_1 the set of morphisms of Q . Let Set/Q_0 denotes the category whose objects are families X of sets X_u indexed by $u \in Q_0$. An element $x \in X_u$ will be called an element over u and we shall sometimes write $d(x)$ for u . Morphisms in Set/Q_0 are families of maps $f_u : X_u \rightarrow Y_u$.

DEFINITION 2.2. Let Q be a quantaloid and X and Y be two objects of Set/Q_0 . A matrix M from X to Y assigns to each pair x, y of $X \times Y$ an element of Q_1 :

$$m_{x,y} : d(x) \rightarrow d(y).$$

Matrices compose by "matrix multiplication": for $M : X \rightarrow Y$, and $N : Y \rightarrow Z$, the composite $M \& N = L : X \rightarrow Z$ has its general element given by

$$l_{x,z} = \bigvee_{y \in Y} m_{x,y} \& n_{y,z}.$$

It is clear that this composition is associative. If Q has units, then the matrix I defined by

$$\begin{aligned} i_{x,x} &= id_{d(x)} \text{ (unit in } Q(d(x), d(x))), \\ i_{x,x'} &= \perp_{d(x), d(x')} \text{ if } x \neq x' \text{ (bottom element in } Q(d(x), d(x'))), \end{aligned}$$

is the unit matrix on X . Matrices from X to Y form a partially ordered set for the point-wise partial order. Sets over Q_0 with matrices as 1-morphisms and 2-morphisms given by partial order determine a bicategory.

3. Q -sets and presheaves on a (midsymmetrical) quantaloid Q

The notions of a (general) Q -set and of a presheaf, which will be given in this section, are taken from [1] and [2].

DEFINITION 3.1. Let Q be a quantaloid. A Q -set is an object X of $Sets/Q_0$ provided with a matrix $A : X \rightarrow X$ satisfying the following:

Idempotency: $A \& A = A$.

An element $x \in X$ of a Q -set (X, A) will said to be strict provided that it satisfies

Strictness: $a_{x,x} \& a_{x,x'} = a_{x,x'}$ and $a_{x',x} \& a_{x,x} = a_{x',x}$ for all $x' \in X$ and a Q -set (X, A) itself will be called strict whenever every element $x \in X$ is strict.

A Q -set (X, A) will be called regular whenever it satisfies

Regularity: $a_{x,x'} = a_{x,x'} \& a_{x',x} \& a_{x,x'}$ for all $x, x' \in X$.

A Q -set (X, A) will be called separated whenever it satisfies

Separation: if $a_{x,x''} = a_{x',x''}$ and $a_{x'',x} = a_{x'',x'}$ for all $x'' \in X$ (which is just the condition that $a_{x,x} = a_{x',x'} = a_{x,x'} = a_{x',x}$ when x and x' are strict), then $x = x'$.

We shall usually write a_x for $a_{x,x}$. Note that regular Q -sets are strict. An example of a regular Q -set is given in Theorem 3.4 (by (1)).

DEFINITION 3.2. Given a strict Q -set (X, A) , the restrictable triplet $(p, x, p^\#)$ (of $Q_1 \times X \times Q_1$) is the element of $Q_1 \times X \times Q_1$ over $d(p, x, p^\#) = \text{dom}(p)$ such that $\text{dom}(p) = \text{cod}(p^\#)$, $\text{cod}(p) = \text{dom}(p^\#) = d(x)$

$$a_x \& p^\# \& p \& a_x \leq a_x, \quad p \& a_x \leq p \& p^\# \& p \& a_x, \quad \text{and} \\ a_x \& p^\# \leq a_x \& p^\# \& p \& p^\#,$$

where the last two inequalities are actually equalities owing to the first inequality.

DEFINITION 3.3. By a *presheaf* on Q will be meant a Q -set (X, A) together with

Restriction: a partial mapping $\uparrow \uparrow: Q_1 \times X \times Q_1 \rightarrow X$ (more precisely, a "matrix" $\uparrow \uparrow = (\uparrow \uparrow_{u,v})_{u,v \in Q_0}^u$ of partial mappings $\uparrow \uparrow_{u,v}: Q(u, v) \times X_v \times Q(v, u) \rightarrow X_u$) from the restrictable triplets $(p, x, p^\#) \in Q_1 \times X \times Q_1$ to the elements $p \uparrow x \uparrow p^\#$ of X satisfying the compatibility conditions that:

$$q \uparrow (p \uparrow x \uparrow p^\#) \uparrow q^\# = (q \& p) \uparrow x \uparrow (p^\# \& q^\#), \quad a_x \uparrow x \uparrow a_x = x, \\ a_{p \uparrow x \uparrow p^\#, x''} = p \& a_{x, x''} \quad \text{and} \quad a_{x'', p \uparrow x \uparrow p^\#} = a_{x'', x} \& p^\#,$$

for all $x, x'' \in X$ and for all restrictable triplets $(p, x, p^\#), (q, p \uparrow x \uparrow p^\#, q^\#) \in Q_1 \times X \times Q_1$. We call a presheaf $(X, A, \uparrow \uparrow)$ on Q regular or separated if the underlying Q -set (X, A) is regular or separated, respectively.

Now we are engaged in an explicit description of the underlying Q -set of a presheaf.

Theorem 3.4. Given a presheaf $(X, A, \uparrow \uparrow)$ on Q , the matrix $A^s = (a_{x,x'}^s)_{x' \in X}^{x \in X}$ defined by

$$a_{x,x'}^s = a_x \& \bigvee \{p \in P_{x,x'} \mid p \& p^\# \uparrow x \uparrow p \& p^\# = p \uparrow x' \uparrow p^\#\} \& a_{x'} \quad (1)$$

for all $x, x' \in X$, satisfies Idempotency and Regularity (and thus makes (X, A^s) into a regular Q -set), where

$$P_{x,x'} = \{p \in Q(d(x), d(x')) \mid \exists p^\# \in Q(d(x'), d(x)) \text{ such that } (p^\#, x, p) \text{ and } \\ (p, x', p^\#) \text{ restrictable (i.e., } a_x \& p \& p^\# \& a_x \leq a_x, \\ p^\# \& a_x = p^\# \& p \& p^\# \& a_x, \quad a_x \& p = a_x \& p \& p^\# \& p, \quad a_{x'} \& p^\# \& p \& a_{x'} \leq a_{x'}, \\ p \& a_{x'} = p \& p^\# \& p \& a_{x'}, \text{ and } a_{x'} \& p^\# = a_{x'} \& p^\# \& p \& p^\#\})\}.$$

Moreover, it is compatible with A and $\uparrow \uparrow$ in the sense that

$$a_{x,x}^s = a_x, \quad (2)$$

$$a_{p \uparrow x \uparrow p^\#, x'}^s = p \& a_{x, x'}^s, \quad \text{and} \quad a_{x', p \uparrow x \uparrow p^\#}^s = a_{x', x} \& p^\#, \quad (3)$$

for all $x, x' \in X$ and all restrictable triplets $(p, x, p^\#) \in Q_1 \times X \times Q_1$ (making the triplet $(X, A^s, \uparrow)$ into a regular presheaf on Q).

Proof. Firstly, note that the restrictability of a triplet $(p^\#, x, p)$ implies the restrictability of the triplet $(p \& p^\#, x, p \& p^\#)$ (used in (1)). For instance, to verify the inequality $a_x \& p \& p^\# \& p \& p^\# \& a_x \leq a_x$, observe that $a_x \& p \& p^\# \& a_x \leq a_x$, in view of midsymmetry, implies that

$$a_x \& p \& p^\# \& p \& p^\# \& a_x = a_x \& p \& p^\# \& a_x \& a_x \& p \& p^\# \& a_x \leq a_x \& a_x = a_x.$$

Verifying each of the axioms of Definition 3.1 for A^s in turn, we argue as follows:

Idempotency: $a_{x,x'}^s = \bigvee_{x'' \in X} a_{x,x''}^s \& a_{x'',x'}^s$, for all $x, x' \in X$, since, on the one hand, we have that

$$\begin{aligned} & \bigvee_{x'' \in X} a_x \& \bigvee \{p \in P_{x,x''} \mid p \& p^\# \uparrow x \uparrow p \& p^\# = p \uparrow x'' \uparrow p^\#\} \& a_{x''} \\ & \& \bigvee \{p' \in P_{x'',x'} \mid p' \& p'^\# \uparrow x'' \uparrow p' \& p'^\# = p' \uparrow x' \uparrow p'^\#\} \& a_{x'} \\ & \leq \bigvee_{x'' \in X} a_x \& \bigvee \{p \& a_{x''} \& p' \in P_{x,x'} \mid (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \uparrow x \uparrow (p \& \\ & a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) = (p \& a_{x''} \& p') \uparrow x' \uparrow (p'^\# \& a_{x''} \& p^\#)\} \& a_{x'} \\ & \leq a_x \bigvee \{p'' \in P_{x,x'} \mid p'' \& p''^\# \uparrow x \uparrow p'' \& p''^\# = p'' \uparrow x' \uparrow p''^\#\} \& a_{x'}, \end{aligned}$$

because, from the midsymmetry of the quantaloid, one deduce that

$$\{p \in P_{x,x''}, p \& p^\# \uparrow x \uparrow p \& p^\# = p \uparrow x'' \uparrow p^\#, p' \in P_{x'',x'}, \text{ and } p' \& p'^\# \uparrow x'' \uparrow p' \& p'^\# = p' \uparrow x' \uparrow p'^\#\} \Rightarrow p \& a_{x''} \& p' \in P_{x,x'}$$

(for instance, to verify the inequality:

$$a_x \& (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \& a_x \leq a_x,$$

observe that relations $a_x \& p \& p^\# \& a_x \leq a_x$, $a_{x''} \& p' \& p'^\# \& a_{x''} \leq a_{x''}$, and $p \& p^\# \& a_x \& p \& p^\# = p \& a_{x''} \& p^\#$, imply that

$$\begin{aligned} a_x \& p \& (a_{x''} \& p' \& p'^\# \& a_{x''}) \& p^\# \& a_x & \leq a_x \& (p \& a_{x''} \& p^\#) \& a_x \\ & = a_x \& p \& p^\# \& a_x \& p \& p^\# \& a_x \leq a_x; \end{aligned}$$

while, for

$$a_x \& (p \& a_{x''} \& p') = a_x \& (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \& (p \& a_{x''} \& p'),$$

we calculate as follows:

$$\begin{aligned} & a_x \& p \& a_{x''} \& (p' \& p'^\#) \& a_{x''} \& (p^\# \& p) \& a_{x''} \& p' \\ & = (a_x \& p \& p^\# \& p) \& (a_{x''} \& p' \& p'^\# \& p') = a_x \& p \& a_{x''} \& p' \end{aligned}$$

and because

$$\begin{aligned}
& \{p \& p^\# \vdash x \frown p \& p^\# = p \vdash x'' \frown p^\# \text{ and } p' \& p'^\# \vdash x'' \frown p' \& p'^\# = p' \vdash x' \frown p'^\#\} \\
& \Rightarrow \{p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \vdash p \& p^\# \vdash x \frown p \& p^\# \frown p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \\
& = p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \vdash p \vdash x'' \frown p^\# \frown p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \text{ and } \\
& p \& a_{x''} \vdash p' \& p'^\# \vdash x'' \frown p' \& p'^\# \frown a_{x''} \& p^\# = p \& a_{x''} \& p' \vdash x' \frown p'^\# \frown a_{x''} \& p^\#\} \\
& \Rightarrow \{p \& a_{x''} \& p' \& p'^\# \& (a_{x''} \& p^\# \& p \& p^\#) \vdash x \frown (p \& p^\# \& p \& a_{x''}) \& p' \& p'^\# \& a_{x''} \& p^\# \\
& = p \& a_{x''} \& (p' \& p'^\#) \& a_{x''} \& (p^\# \& p) \& a_{x''} \vdash x'' \frown a_{x''} \& (p^\# \& p) \& a_{x''} \& (p' \& p'^\#) \& a_{x''} \& p^\# \\
& \text{and } p \& a_{x''} \& p' \& p'^\# \vdash x'' \frown p' \& p'^\# \& a_{x''} \& p^\# = p \& a_{x''} \& p' \vdash x' \frown p'^\# \& a_{x''} \& p^\#\} \\
& \Rightarrow \{p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \vdash x \frown p \& a_{x''} \& p' \& p'^\# \& a_{x''} \& p^\# \\
& = (p \& p^\# \& p \& a_{x''}) \& p' \& p'^\# \& a_{x''} \vdash x'' \frown a_{x''} \& p' \& p'^\# \& (a_{x''} \& p^\# \& p \& p^\#) \\
& \text{and } p \& a_{x''} \& p' \& p'^\# \vdash x'' \frown p' \& p'^\# \& a_{x''} \& p^\# = p \& p' \vdash x' \frown p'^\# \& a_{x''} \& p^\#\} \\
& \Rightarrow (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \vdash x \frown (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \\
& = p \& a_{x''} \& p' \& p'^\# \vdash x'' \frown p' \& p'^\# \& a_{x''} \& p^\# = (p \& a_{x''} \& p') \vdash x' \frown (p'^\# \& a_{x''} \& p^\#) \\
& \Rightarrow (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \vdash x \frown (p \& a_{x''} \& p') \& (p'^\# \& a_{x''} \& p^\#) \\
& = (p \& a_{x''} \& p') \vdash x' \frown (p'^\# \& a_{x''} \& p^\#),
\end{aligned}$$

while, on the other hand,

$$\begin{aligned}
& \bigvee_{x'' \in X} a_x \& \bigvee \{p \in P_{x, x''} | p \& p^\# \vdash x \frown p \& p^\# = p \vdash x'' \frown p^\#\} \& a_{x''} \\
& \& a_{x''} \& \bigvee \{p' \in P_{x'', x'} | p' \& p'^\# \vdash x'' \frown p' \& p'^\# = p' \vdash x' \frown p'^\#\} \& a_{x'} \\
& \geq a_x \& \bigvee \{p \in P_{x, x'} | p \& p^\# \vdash x \frown p \& p^\# = p \vdash x' \frown p^\#\} \& a_{x'}
\end{aligned}$$

(putting $x'' = x'$ and $p' = a_{x'}$);

Regularity: $a_{x, x'}^s = a_{x, x}^s \& a_{x', x}^s \& a_{x, x'}^s$, for all $x, x' \in X$, since, in the one direction, we have that

$$a_{x, x'}^s \& a_{x', x}^s \& a_{x, x'}^s \leq a_{x, x'}^s \text{ (by Idempotency of } A^s \text{ twice),}$$

while, in the other,

$$\begin{aligned}
& a_x \& \bigvee \{p \in P_{x, x'} | p \& p^\# \vdash x \frown p \& p^\# = p \vdash x' \frown p^\#\} \& a_{x'} \\
& \& a_{x'} \& \bigvee \{q \in P_{x', x} | q \& q^\# \vdash x' \frown q \& q^\# = q \vdash x \frown q^\#\} \& a_x \\
& \& a_{x'} \& \bigvee \{r \in P_{x, x'} | r \& r^\# \vdash x \frown r \& r^\# = r \vdash x' \frown r^\#\} \& a_{x'} \\
& \geq a_x \& \bigvee \{p \& p^\# \& p \& a_{x'} (= p \& a_{x'}) | p \in P_{x, x'}, p \& p^\# \vdash x \frown p \& p^\# = p \vdash x' \frown p^\#\} \& a_{x'}
\end{aligned}$$

(putting $q = r^\# = p^\#$ and $q^\# = r = p$ and noting that

$$p^\# \& p \vdash x' \frown p^\# \& p = p^\# \vdash x \frown p \Leftrightarrow p \& p^\# \vdash x \frown p \& p^\# = p \vdash x' \frown p^\#$$

for all $p \in P_{x,x'}$, which completes the verification that the underlying family X of sets of the presheaf $(X, A, \dashv)$ together with A^s is a Q -set.

For (2), in the one direction, we have that

$$a_x = a_x \& a_x \& a_x \leq a_x \& \bigvee \{p \in P_{x,x} \mid p \& p^\# \dashv x \dashv p \& p^\# = p \dashv x' \dashv p^\#\} \& a_x,$$

while, in the other,

$$a_x \& \bigvee \{p \in P_{x,x} \mid p \& p^\# \dashv x \dashv p \& p^\# = p \dashv x' \dashv p^\#\} \& a_x \leq a_x,$$

since from the relation $p \& p^\# \dashv x \dashv p \& p^\# = p \dashv x \dashv p^\#$, it follows that

$$\begin{aligned} a_x \& p \& a_x &= a_x \& a_{p \dashv x \dashv p^\#, x} = a_x \& a_{p \& p^\# \dashv x \dashv p \& p^\#, x} \\ &= a_x \& p \& p^\# \& a_x \leq a_x \quad (\text{by the restrictability of } (p, x, p^\#)). \end{aligned}$$

To prove the first relation in (3), by the observation that

$$q \in P_{p \dashv x \dashv p^\#, x'} \Rightarrow a_x \& p^\# \& q \& a_{x'} \in P_{x, x'}$$

and that

$$\begin{aligned} q \& q^\# \dashv p \dashv x \dashv p^\# \dashv q \& q^\# &= q \dashv x' \dashv q^\# \\ \Rightarrow a_x \& p^\# \& q \& a_{x'} \& q^\# \dashv p \dashv x \dashv p^\# \dashv q \& q^\# \dashv q \& a_{x'} \& q^\# \& p \& a_x \\ &= a_x \& p^\# \& q \& a_{x'} \& q^\# \dashv q \dashv x' \dashv q^\# \dashv q \& a_{x'} \& q^\# \& p \& a_x \\ \Rightarrow a_x \& p^\# \& q \& (a_{x'} \& q^\# \& q \& q^\#) \& p \dashv x \dashv p^\# \& (q \& q^\# \& q \& a_{x'}) \& q^\# \& p \& a_x \\ &= a_x \& p^\# \& (q \& a_{x'} \& q^\# \& q \& a_{x'}) \dashv x' \dashv (a_{x'} \& q^\# \& q \& a_{x'} \& q^\#) \& p \& a_x \\ \Rightarrow (a_x \& p^\# \& q \& a_{x'}) \& (a_{x'} \& q^\# \& p \& a_x) \dashv x \dashv (a_x \& p^\# \& q \& a_{x'}) \\ &\quad \& (a_{x'} \& q^\# \& p \& a_x) = (a_x \& p^\# \& q \& a_{x'}) \dashv x' \dashv (a_{x'} \& q^\# \& p \& a_x), \end{aligned}$$

we obtain the following estimations

$$\begin{aligned} a_{p \dashv x \dashv p^\#, x'}^s &= a_{p \dashv x \dashv p^\#, x'} \& \bigvee \{q \in P_{p \dashv x \dashv p^\#, x'} \mid q \& q^\# \dashv p \dashv x \dashv p^\# \dashv q \& q^\# \\ &= q \dashv x' \dashv q^\#\} \& a_{x'} \leq p \& \bigvee \{a_x \& p^\# \& q \& a_{x'} \in P_{x, x'} \mid (a_x \& p^\# \& q \& a_{x'}) \\ &\quad \& (a_{x'} \& q^\# \& p \& a_x) \dashv x \dashv (a_x \& p^\# \& q \& a_{x'}) \& (a_{x'} \& q^\# \& p \& a_x) \\ &= (a_x \& p^\# \& q \& a_{x'}) \dashv x' \dashv (a_{x'} \& q^\# \& p \& a_x)\} \\ &\leq p \& \bigvee \{q' \in P_{x, x'} \mid q' \& q'^\# \dashv x \dashv q' \& q'^\# = q' \dashv x' \dashv q'^\#\}, \end{aligned}$$

which implies that

$$a_{p \dashv x \dashv p^\#, x'}^s \leq p \& a_{x, x'}^s.$$

For the converse, using the implication

$$\{(p, x, p^\#) \text{ is restrictable, } r \in P_{x, x'}, \text{ and } r \& r^\# \vdash x \uparrow r \& r^\# = r \vdash x' \uparrow r^\#\} \\ \text{(from which one has } r \& a_{x'} = r \& r^\# \& a_{x, x'})\} \Rightarrow p \& a_x \& r \in P_{p \vdash x \uparrow p^\#, x'}$$

and the following series of implications

$$\begin{aligned} & \{(p, x, p^\#) \text{ is restrictable and } r \& r^\# \vdash x \uparrow r \& r^\# = r \vdash x' \uparrow r^\#\} \\ & \Rightarrow p \& a_x \& r \& r^\# \vdash r \& r^\# \vdash x \uparrow r \& r^\# \uparrow r \& r^\# \& a_x \& p^\# \\ & = p \& a_x \& r \& r^\# \vdash r \vdash x' \uparrow r^\# \uparrow r \& r^\# \& a_x \& p^\# \\ & \Rightarrow p \& (a_x \& r \& r^\# \& r) \& r^\# \vdash x \uparrow r \& (r^\# \& r \& r^\# \& a_x) \& p^\# \\ & = p \& (a_x \& r \& r^\# \& r) \vdash x' \uparrow (r^\# \& r \& r^\# \& a_x) \& p^\# \\ & \Rightarrow p \& (p^\# \& p) \& a_x \& (r \& r^\#) \& a_x \vdash x \uparrow a_x \& (r \& r^\#) \\ & \quad \& a_x \& (p^\# \& p) \& p^\# = (p \& a_x \& r) \vdash x' \uparrow (r^\# \& a_x \& p^\#) \\ & \Rightarrow (p \& a_x \& r) \& (r^\# \& a_x \& p^\#) \& (p \& a_x \vdash x \uparrow a_x \& p^\#) \\ & \& (p \& a_x \& r) \& (r^\# \& a_x \& p^\#) = (p \& a_x \& r) \vdash x' \uparrow (r^\# \& a_x \& p^\#) \\ & \Rightarrow (p \& a_x \& r) \& (r^\# \& a_x \& p^\#) \vdash (p \vdash x \uparrow p^\#) \uparrow (p \& a_x \& r) \\ & \quad \& (r^\# \& a_x \& p^\#) = (p \& a_x \& r) \vdash x' \uparrow (r^\# \& a_x \& p^\#), \end{aligned}$$

we calculate

$$\begin{aligned} p \& a_{x, x'}^s &= p \& a_x \& \bigvee \{r \in P_{x, x'} \mid r \& r^\# \vdash x \uparrow r \& r^\# = r \vdash x' \uparrow r^\#\} \& a_{x'} \\ &= p \& (p^\# \& p) \& a_x \& \bigvee \{a_x \& r \mid r \in P_{x, x'}, r \& r^\# \vdash x \uparrow r \& r^\# = r \vdash x' \uparrow r^\#\} \& a_{x'} \\ &= (p \& a_x \& p^\#) \& \bigvee \{p \& a_x \& r \mid r \in P_{x, x'}, r \& r^\# \vdash x \uparrow r \& r^\# = r \vdash x' \uparrow r^\#\} \& a_{x'} \\ &\leq a_{p \vdash x \uparrow p^\#, x'} \& \bigvee \{p \& a_x \& r \in P_{p \vdash x \uparrow p^\#, x'} \mid (p \& a_x \& r) \& (r^\# \& a_x \& p^\#) \vdash (p \vdash x \uparrow p^\#) \uparrow \\ & \quad (p \& a_x \& r) \& (r^\# \& a_x \& p^\#) = (p \& a_x \& r) \vdash x' \uparrow (r^\# \& a_x \& p^\#)\} \& a_{x'} \\ &\leq a_{p \vdash x \uparrow p^\#, x'} \& \bigvee \{r' \in P_{p \vdash x \uparrow p^\#, x'} \mid r' \& r'^\# \vdash (p \vdash x \uparrow p^\#) \uparrow r' \& r'^\# = r' \vdash x' \uparrow r'^\#\} \& a_{x'} \\ &= a_{p \vdash x \uparrow p^\#, x'}^s, \end{aligned}$$

which completes the proof of the first relation in (3), while the verification of the rest is similar.

We obtain further properties of the regular Q -set (X, A^s) (introduced in Theorem 3.4) in the next few propositions.

Lemma 3.5. *Let $(X, A, \vdash \uparrow)$ be a presheaf on Q and $A^s = (a_{x, x'}^s)_{x' \in X}^{x \in X}$ be the structural matrix defined by (1) (in Theorem 3.4). Then the inequality*

$$a_{x, x'}^s \leq a_{x, x'} \quad (4)$$

holds for all $x, x' \in X$, i.e., $A^s \leq A$.

Proof. Let us consider elements $x, x' \in X$ and $p, p^\# \in Q_1$ with the properties that $(p^\#, x, p)$ and $(p, x', p^\#)$ are restrictable (i.e., $p \in P_{x, x'}$) and that the equality $p \& p^\# \upharpoonright x \upharpoonright p \& p^\# = p \upharpoonright x' \upharpoonright p^\#$ holds. Then we infer from (3) that $p \& p^\# \& a_{x, x'} = p \& a_{x'}$. Now we obtain

$$a_x \& p \& a_{x'} = a_x \& p \& p^\# \& a_{x, x'} = (a_x \& p \& p^\# \& a_x) \& a_{x, x'} \leq a_x \& a_{x, x'} = a_{x, x'},$$

whence (4).

Lemma 3.6. *Let $(X, A, \upharpoonright)$ be a separated presheaf on Q (i.e., (X, A) be a separated Q -set). Let $(p, x', p^\#)$ and $(p^\#, x, p)$ be restrictable triplets. Then the relation*

$$p \& p^\# \upharpoonright x \upharpoonright p \& p^\# = p \upharpoonright x' \upharpoonright p^\# \quad (5)$$

(presented in (1)) is just the condition that

$$p \& p^\# \& a_x \& p \& p^\# = p \& a_{x'} \& p^\# = p \& p^\# \& a_{x, x'} \& p^\# = p \& a_{x', x} \& p \& p^\#. \quad (6)$$

Proof. To prove the implication from (5) to (6), assume (5). Then

$$a_{p \& p^\# \upharpoonright x \upharpoonright p \& p^\#} = a_{p \upharpoonright x' \upharpoonright p^\#} = a_{p \& p^\# \upharpoonright x \upharpoonright p \& p^\#, p \upharpoonright x' \upharpoonright p^\#} = a_{p \upharpoonright x' \upharpoonright p^\#, p \& p^\# \upharpoonright x \upharpoonright p \& p^\#}, \quad (7)$$

which implies (6). Assume now (6). Then it is clear that the relations (7) hold. By Separation, this means that (5) holds.

Lemma 3.7. *If a presheaf $(X, A, \upharpoonright)$ is separated, then the triplets $(a_{x', x}, x, a_{x, x'})$ and $(a_{x, x'}, x', a_{x', x})$ are restrictable (i.e., $a_{x, x'} \in P_{x, x'}$) and*

$$a_{x, x'} \& a_{x', x} \upharpoonright x \upharpoonright a_{x, x'} \& a_{x', x} = a_{x, x'} \upharpoonright x' \upharpoonright a_{x', x}$$

for all $x, x' \in X$.

Proof. In view of Idempotency and Regularity of A , it is clear that $a_{x, x'} \in P_{x, x'}$ for all $x, x' \in X$. Moreover, we have that

$$a_{x, x'} \& a_{x', x} \& a_x \& a_{x, x'} \& a_{x', x} = a_{x, x'} \& a_{x'} \& a_{x', x} = a_{x, x'} \& a_{x', x} \& a_{x, x'} \& a_{x', x}$$

for all $x, x' \in X$. By Lemma 3.6, this just says the statement is correct.

Theorem 3.8. *If $(X, A, \upharpoonright)$ is a regular and separated presheaf, then the Q -set (X, A^s) introduced in Theorem 3.4 (by (1)) is exactly the underlying Q -set (X, A) , that is,*

$$a_{x, x'} = a_x \& \bigvee \{p \in P_{x, x'} \mid p \& p^\# \upharpoonright x \upharpoonright p \& p^\# = p \upharpoonright x' \upharpoonright p^\#\} \& a_{x'}$$

for all $x, x' \in X$.

Proof. From Lemma 3.7, it follows that

$$a_{x,x'} = a_x \& a_{x,x'} \& a_{x'} \leq a_x \& \bigvee \{p \in P_{x,x'} \mid p \& x p^\# \uparrow x \uparrow p \& p^\# = p \uparrow x' \uparrow p^\#\} \& a_{x'}$$

for all $x, x' \in X$. In view of Lemma 3.5, this concludes the proof.

References

- [1] G. Van den Bossche, Quantaloids and non-commutative ring representations, *Appl. Categ. Structures*, **3**, 305–320 (1995).
- [2] R.P. Gylys, Sheaves on quantaloids, *Liet. matem. rink.*, **40**(2), 133–171 (2000).
- [3] C.J. Mulvey, M. Nawaz, Quantaes: quantal sets, In *Non-Classical Logics and their Applications to Fuzzy Subsets*, eds. U. Höhle, E.P. Klement, Kluwer, Boston, 159–217 (1995).

Priešpluoštai virš midsimetriškų kvantaloidų

R.P. Gylys

Ankstesniame straipsnyje [2] mes nagrinėjome priešpluoščius ir pluoštus virš laisvojo kvantaloido. Šiame darbe tiriami priešpluoščiai virš midsimetriško kvantaloido. Išvesta priešpluoščio struktūrinės matricos išraiška per jo siaurinę ir diagonalę.